

Fig. 1

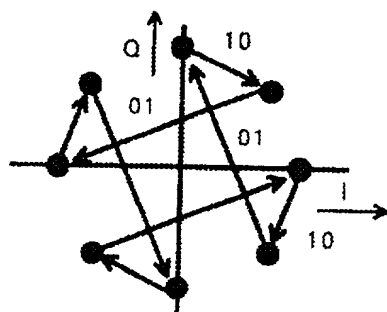


Fig. 2

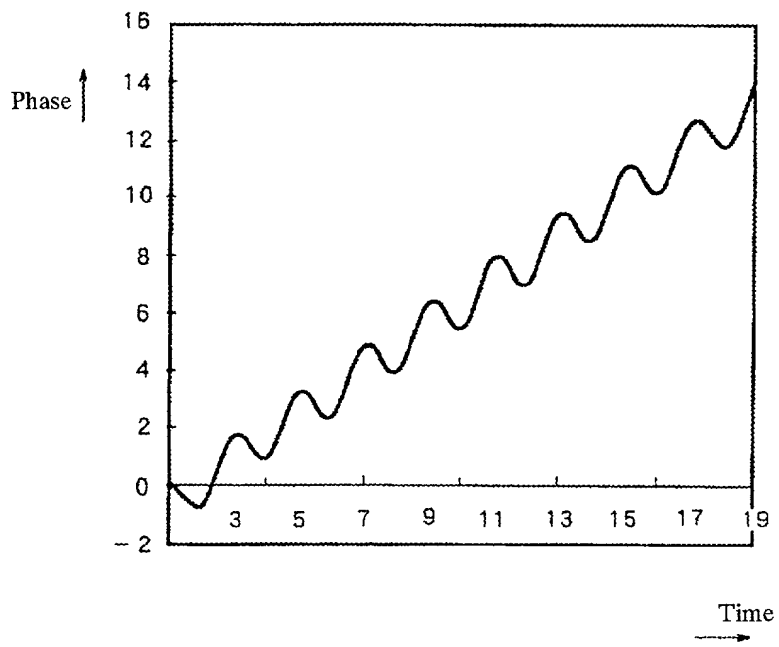


Fig. 3

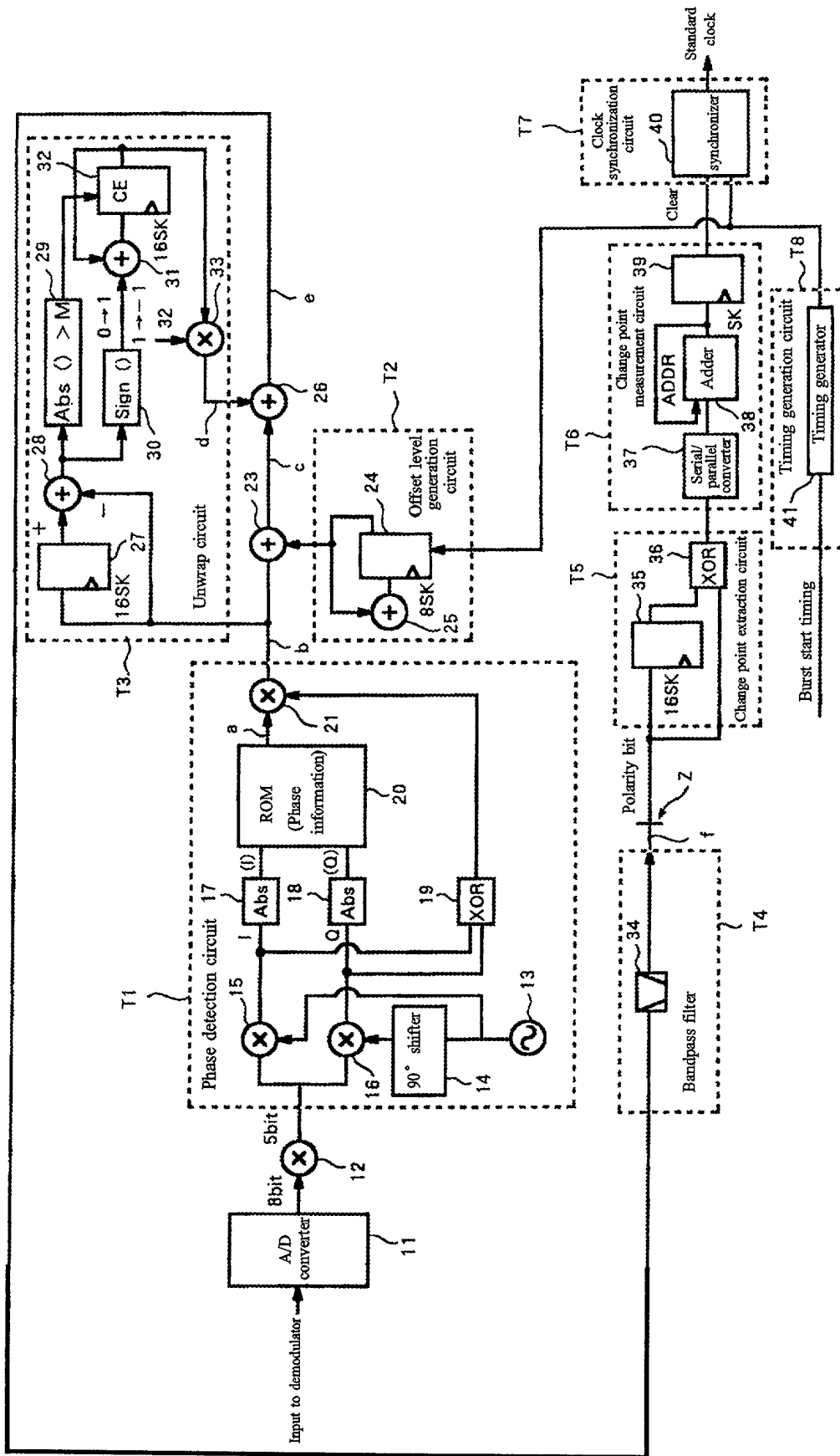
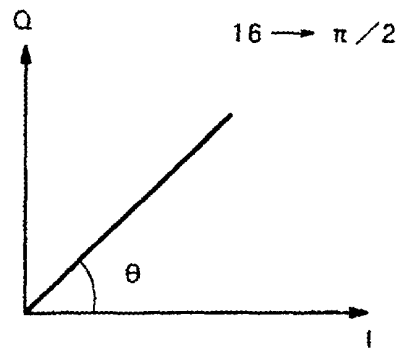


Fig. 4

(a)

$Q \backslash I$	0000	0001	0010	...
0000	a1	a2	a3	...
0001	a4	a5	a6	...
0010	a7	a8	a9	...
...

(b)



(c)

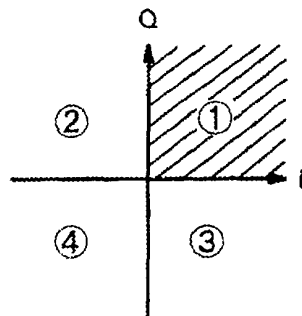


Fig. 5

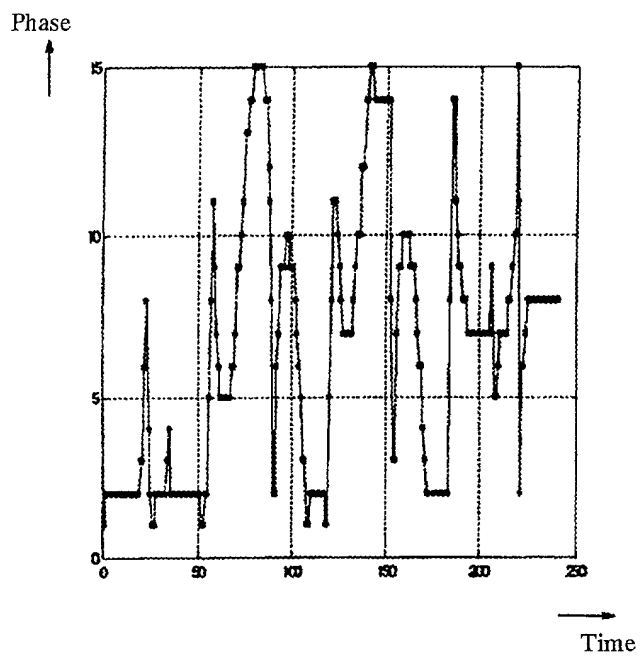


Fig. 6

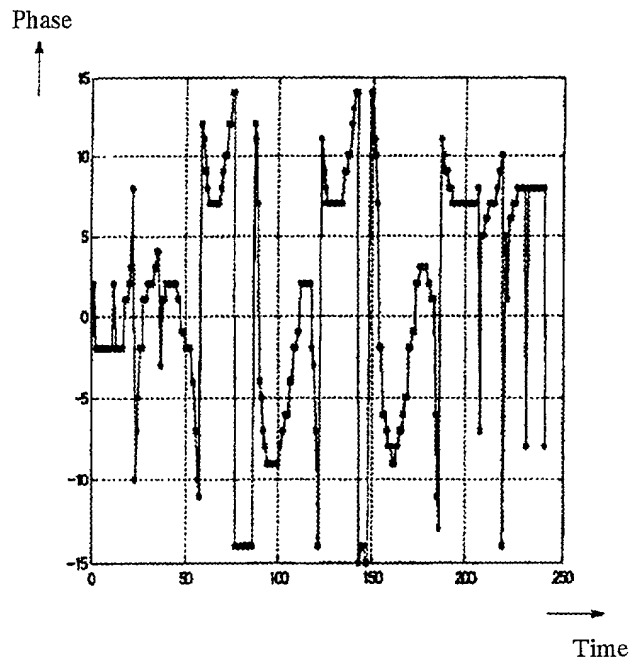
[illegible]

Fig. 7

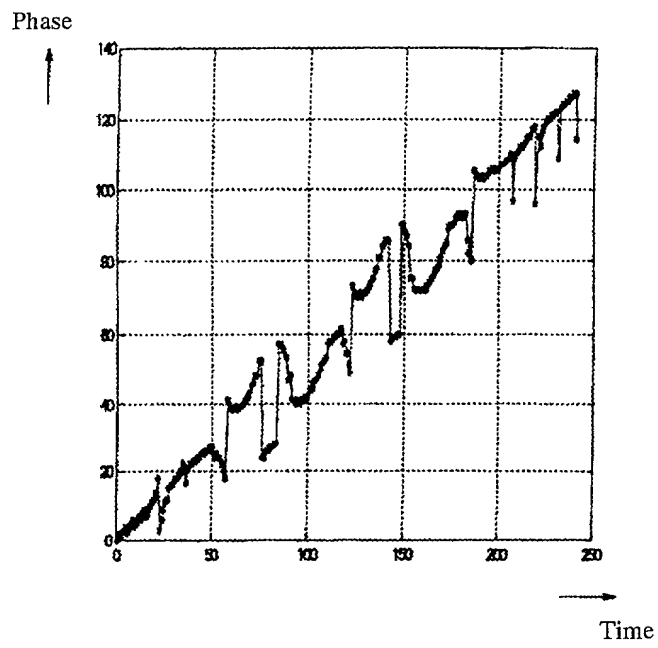


Fig. 8

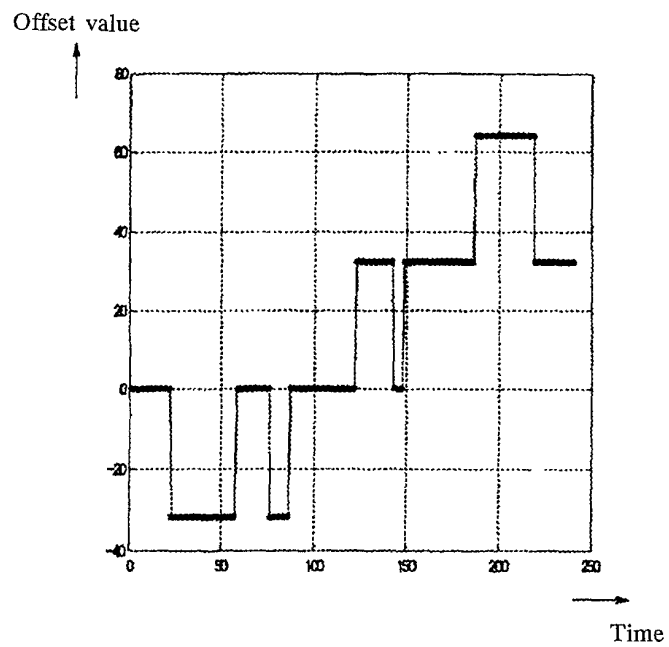


Fig. 9

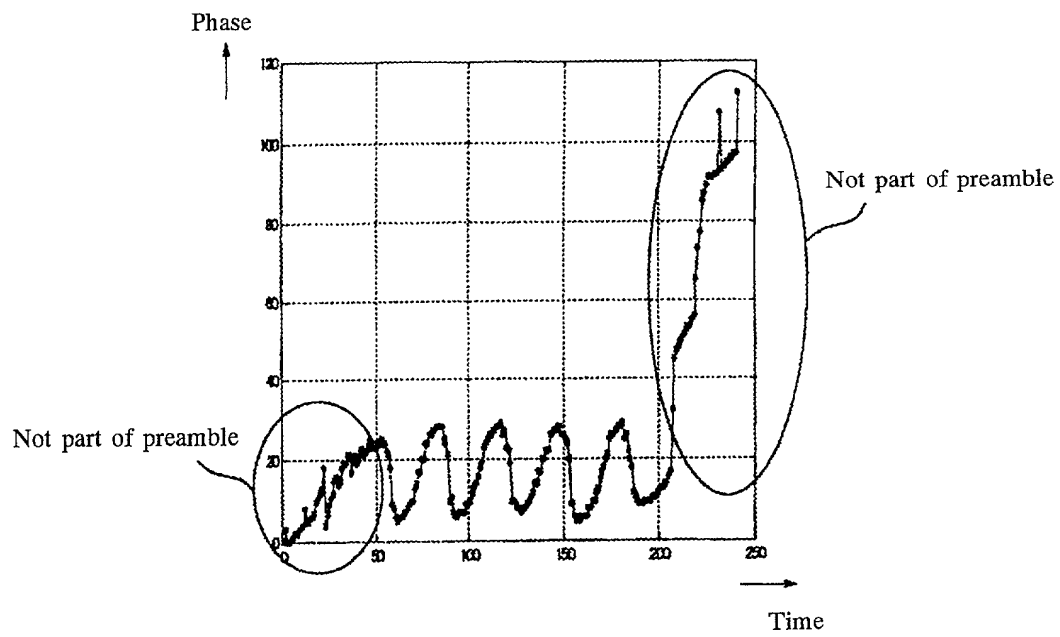


Fig. 10

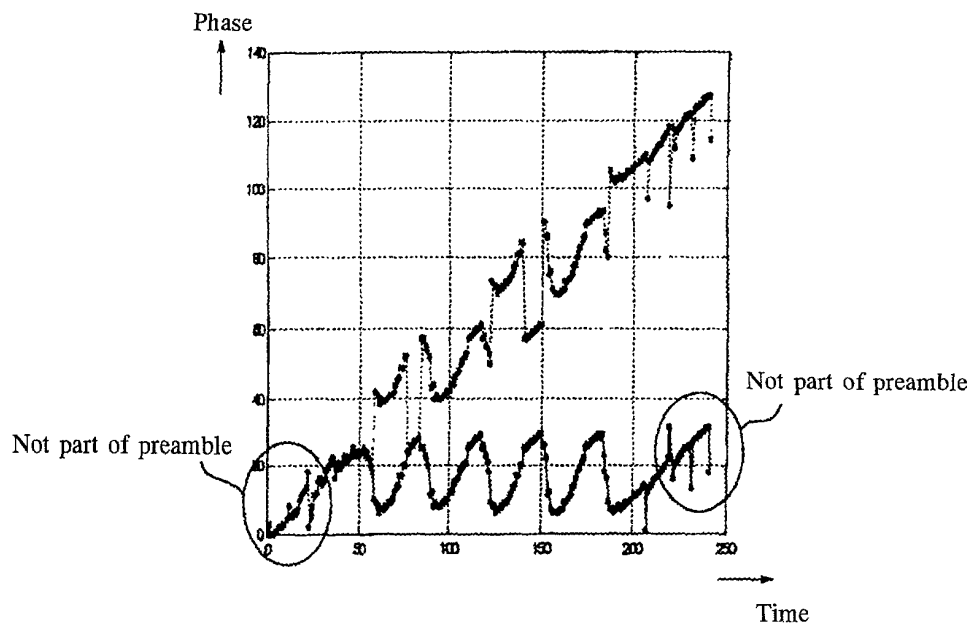


Fig. 11

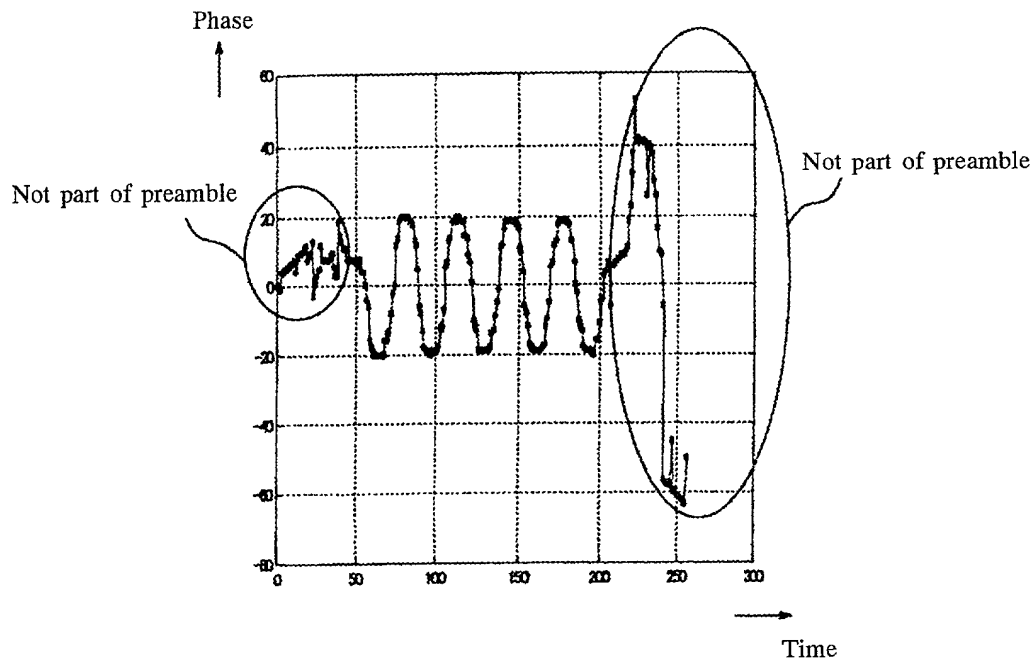


Fig. 12

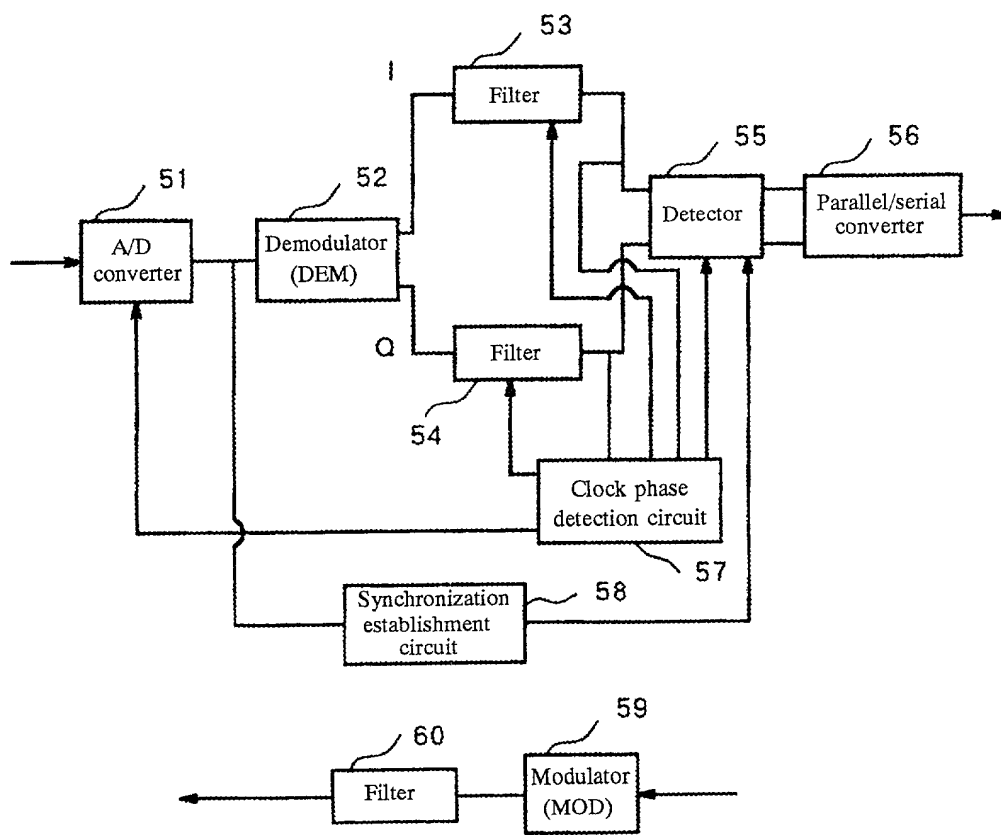


Fig. 13

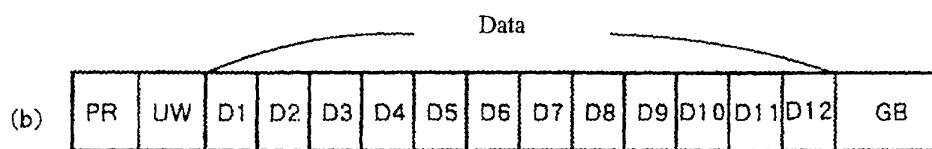
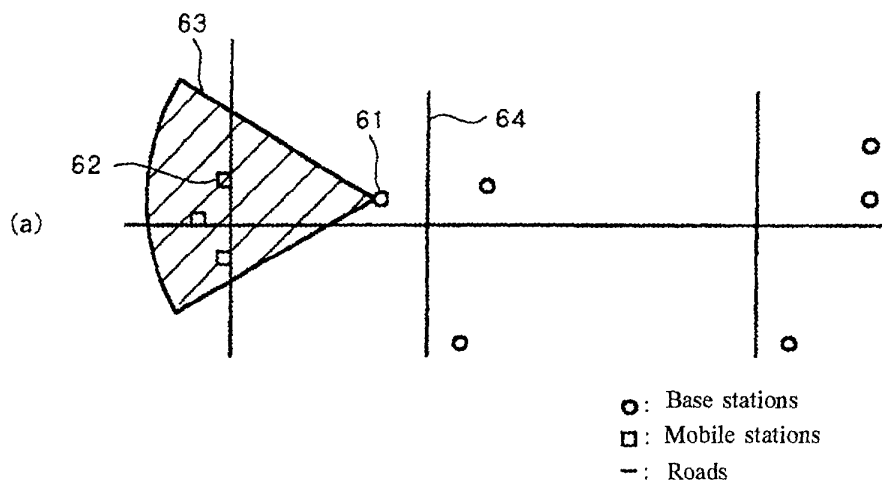


Fig. 14

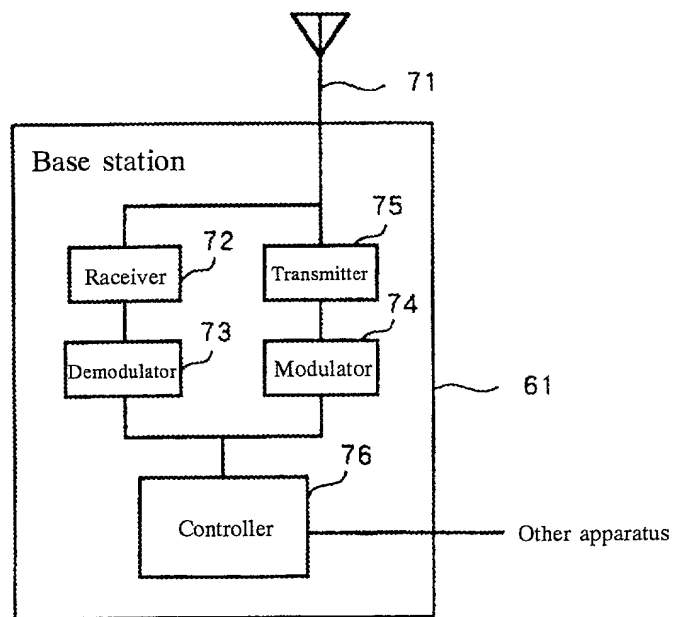


Fig. 15

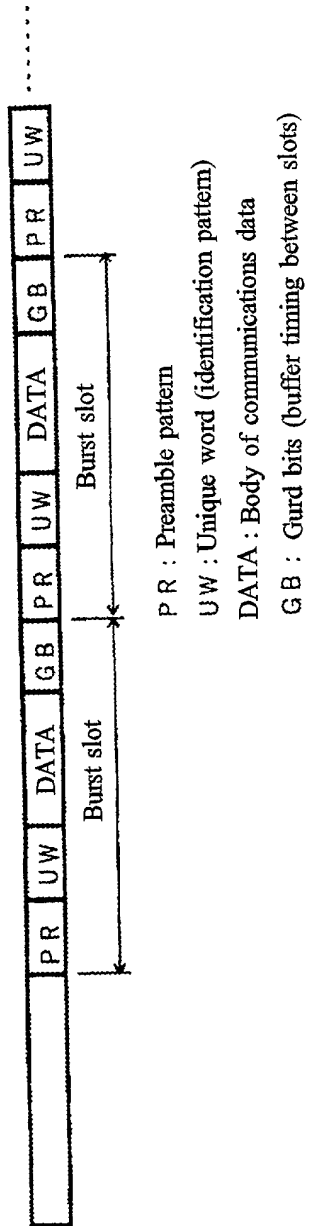


Fig. 16

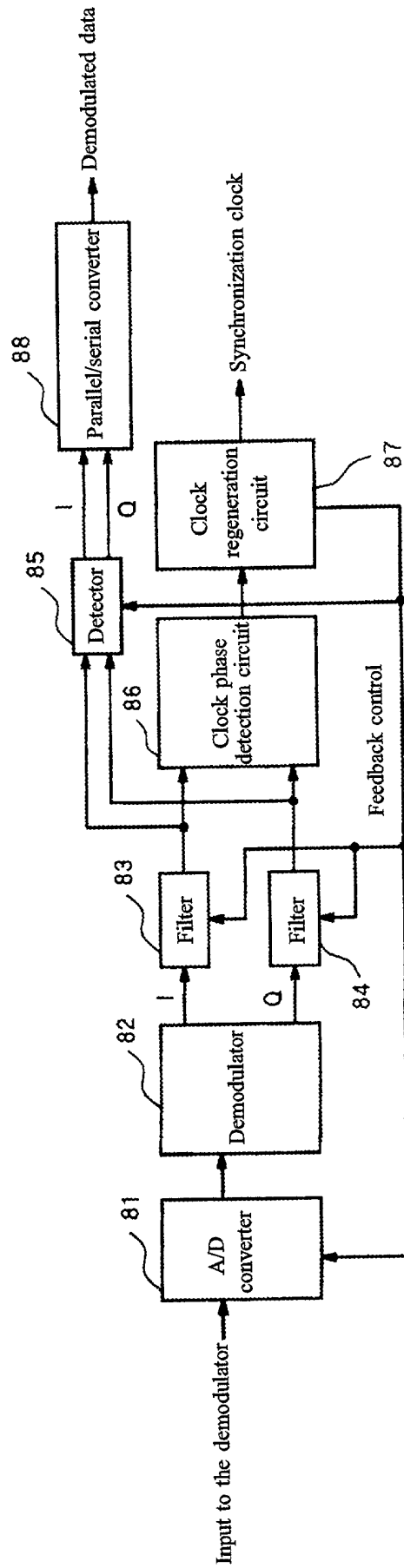
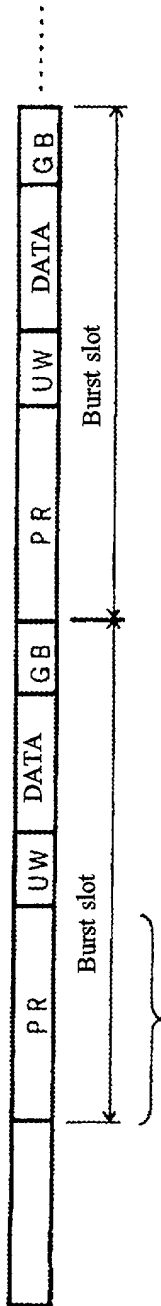


Fig. 17

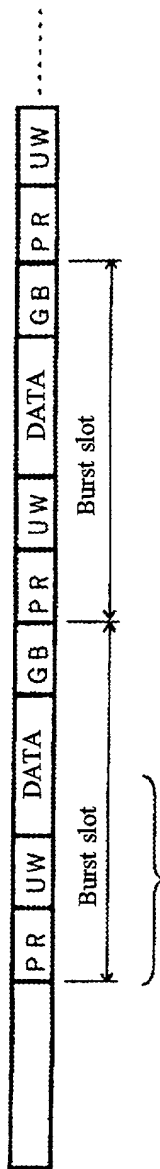
(a) Case in which preamble is lengthened



Preamble period is made 100 symbols or longer

In this case, the length of the preamble becomes a large fraction of the entire length, so the data transmission rate drops.

(b) Case in which first burst is discarded



If the preamble is not made long, there is a risk that data cannot be demodulated with only a length of roughly 100 symbols. In the above diagram, there is a risk that the portion up to UW and DATA cannot be received correctly at the time of receipt of the first burst. In addition, it is necessary to hold the timing of synchronization at the time of receipt of the first burst so that it can be used in the second and subsequent bursts.

PR : Preamble pattern

UW : Unique word (identification pattern)

DATA : Body of communications data

GB : Guard bits (buffer timing between slots)

Fig. 18